# Course 18.312: Algebraic Combinatorics 

Lecture Notes \# 23-24 Addendum by Gregg Musiker
April 6th - 8th, 2009

The following is an outline of the material covered April 6th and 8th in class. This material can be found in Chapter 5 of Stanley's Enumerative Combinatorics Volume 2. Proofs of most of the results are in class notes.

## 1 Exponential Generating Functions

Definition. Given $f, g: \mathbb{N} \rightarrow \mathbb{Z}$, which we think of as counting objects of sizes $k$ in two set $\mathcal{F}$ and $\mathcal{G}$, respectively, we define a new function $h: \mathbb{N} \rightarrow \mathbb{Z}$ by the following:

$$
h(\# X)=\sum_{(S, T)} f(\# S) g(\# T)
$$

where $X$ is a finite set and $(S, T)$ disjointly partition $X$, i.e. $S \cap T=\emptyset$ and $S \cup T=X$. Sets $S$ and $T$ are allowed to be empty.

Definition. We define the exponential generating function of sequence $\{f(n)\}$ to be

$$
E_{f}(x):=\sum_{n \geq 0} f(n) \frac{x^{n}}{n!}
$$

## Proposition.

$$
E_{h}(x)=E_{f}(x) E_{g}(x) .
$$

The following is Corollary 5.1.6 of Stanley's Enumerative Combinatorics 2.
Theorem. (The Exponential Formula) Given $f:\{1,2, \ldots\} \rightarrow \mathbb{Z}$, define a new function $h: \mathbb{N} \rightarrow \mathbb{Z}$ by $h(0)=1$ and

$$
h(\# S)=\sum_{k \geq 1} \sum_{B_{1}, \ldots, B_{k}} f\left(\# B_{1}\right) f\left(\# B_{2}\right) \cdots f\left(\# B_{k}\right)
$$

for $\# S \geq 1$. Here, the sum is over partitions of $S$, i.e. $B_{i} \cap B_{j}=\emptyset$ for all $i \neq j$. We assume these blocks $B_{i}$ are non-empty, and $B_{1} \cup B_{2} \cup \cdots \cup B_{k}=S$. Then

$$
E_{h}(x)=\exp \left(E_{f}(x)\right) .
$$

The following is Corollary 5.1.8 of Stanley's Enumerative Combinatorics 2.
Theorem. (Permutation Version of the Exponential Formula) Given $f:\{1,2, \ldots\} \rightarrow$ $\mathbb{Z}$, define a new function $h: \mathbb{N} \rightarrow \mathbb{Z}$ by $h(0)=1$ and let $n=\# S$,

$$
h(n)=\sum_{\pi \in S_{n}} f\left(\# C_{1}\right) f\left(\# C_{2}\right) \cdots f\left(\# C_{k}\right)
$$

for $\# S \geq 1$. Here, the $C_{i}$ 's are the cycles, thought of as sets of $S$, in the disjoint cycle decmposition of $\pi$. Then

$$
E_{h}(x)=\exp \left(\sum_{n \geq 1} f(n) \frac{x^{n}}{n}\right)
$$

Application: The nubmer of simple graphs on $n$ vertices is $2\binom{n}{2}$ and we let $c(n)$ be the number of connected graphs on $n$ vertices.

$$
\exp \left(\sum_{n \geq 1} c(n) \frac{x^{n}}{n!}\right)=\sum_{n \geq 0} 2^{\binom{n}{2}} \frac{x^{n}}{n!}
$$

## 2 Tree Enumeration

A tree is an undirected graph with no cycles. A tree is rooted if it has a distinguished vertex (called the root).

Let $T_{n}=\#$ labeled trees on $n$ vertices.

Let $t_{n}=\#$ labeled rooted trees on $n$ vertices.
A forest is a disjoint union of trees. A rooted forest is a collection of rooted trees, one root for each tree.

Let $f_{n}=\#$ of rooted labeled forests on $n$ vertices.
Claim: $T_{n+1}=f_{n}$ and $t_{n}=n T_{n}$.

Bijective Proofs: Peel off root, labeled $(n+1)$ of a rooted tree and left with a rooted forest. A rooted tree is a choice of a labeled tree plus a choice of a vertex to be the root.

A Rooted Forest is a collection of rooted trees, so we can use the exponential formula to count. Let

$$
\begin{gathered}
y=E_{t}(x)=\sum_{n \geq 1} t_{n} \frac{x^{n}}{n!} \text { and } E_{f}(x)=\sum_{n \geq 0} f_{n} \frac{x^{n}}{n!} . \\
E_{f}(x)=\exp (y) . \text { On the other hand, } t_{n+1}=(n+1) f_{n} \text {, so } \\
x E_{f}(x)=\sum_{n \geq 0} f_{n} \frac{x^{n+1}}{n!}=\sum_{n \geq 0} t_{n+1} \frac{x^{n+1}}{(n+1)!}=E_{t}(x)=y
\end{gathered}
$$

Thus $y=E_{t}(x)$ satisfies $x e^{y}=y$. We can solve this identity in a way that allows us to compute coefficients of $y$ using a technique known as the Lagrange Inversion Formula.

But first, we compute $t_{n}$ 's combinatorially:
Claim. There are $\binom{n}{d_{1}, d_{2}, \ldots, d_{n}}=\frac{(n-1)!}{d_{1}!l_{2}!\cdots d_{n}!}$ rooted trees on $\{1,2, \ldots, n\}$ in which vertex $i$ has outdegree $d_{i}$, where the outdegree of a vertex $v_{i}$ is the number of its neighbors further away from the root. These neighbors are called children and the unique neighbor closer to the root is called a parent. A vertex with no children is called a leaf. (Notice that $\sum_{i=1}^{n} d_{i}=n-1$.)

We prove this claim using the Prüfer code. Start with a rooted labeled tree $T$.

1. Locate the leaf with the smallest label.
2. Write down the label of its unique parent. Delete this leaf and its adjoing edges.
3. Go to step 1.

Application: The Prüfer code gives bijections between desired set of sequences and rooted trees with specified outdegrees.

Corollary. $t(n)=n^{n-1}$, the number of sequences of length $(n-1)$ on $n$ letters.
Corollary (Cayley's Theorem). $T(n)=n^{n-2}$, the number of (unrooted) labeled trees on $n$ vertices.

Remark. The Catalan numbers count binary trees in several different ways.

## 3 Statement of Lagrange Inversion

Given a formal power series $f(x)=a_{1} x+a_{2} x^{2}+a_{3} x^{3}+\ldots$, we say that $f(x)$ has a compositional inverse $f^{\langle-1\rangle}(x)=g(x)=b_{1} x+b_{2} x^{2}+b_{3} x^{3}+\ldots$ if $f(g(x))=$ $g(f(x))=x$.

Proposition. $f(x)$ has a compositional inverse iff $a_{1} \neq 0$. In this case, the compositional inverse is unique.

Note that
$a_{1}\left(b_{1} x+b_{2} x^{2}+b_{3} x^{3}+\ldots\right)+a_{2}\left(b_{1} x+b_{2} x^{2}+\ldots\right)^{2}+a_{3}\left(b_{1} x+\ldots\right)^{3}+\cdots=x+0 x^{2}+0 x^{3}+\ldots$
if and only if

$$
\begin{aligned}
a_{1} b_{1} & =1 \\
a_{1} b_{2}+a_{2} b_{1}^{2} & =0 \\
a_{1} b_{3}+2 a_{2} b_{1} b_{2}+a_{3} b_{1}^{3} & =0
\end{aligned}
$$

Theorem (Lagrange Inversion Formula). In particular,

$$
\left[x^{n}\right] F^{\langle-1\rangle}(x)=\frac{1}{n}\left[x^{n-1}\right]\left(\frac{x}{F(x)}\right)^{n}
$$

where the right-hand-side can be written equivalently as $\frac{1}{n}\left[x^{-1}\right] F(x)^{-n}$.
Exercise 1: Let $F(x)=\sum_{k \geq 1} \frac{x^{k}}{k!}$ and show that $F^{\langle-1\rangle}(x)=\sum_{k \geq 1} \frac{(-1)^{k+1}}{k} x^{k}$.
(Hint: You will also recognize these as power series of familiar functions.
Exercise 2: Let $F(x)=x e^{-x}$ and we have $E_{t}(x)=F^{\langle-1\rangle}(x)$. Also

$$
\frac{1}{n}\left[x^{n-1}\right]\left(\frac{x}{x e^{-x}}\right)^{n}=\frac{1}{n}\left[x^{n-1}\right] e^{n x}=\frac{1}{n} \frac{n^{n-1}}{(n-1)!}=\frac{n^{n-1}}{n!} .
$$

Consequently, we obtain a second proof that $t_{n}=n^{n-1}$.

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