

1.225J (ESD 205) Transportation Flow Systems

Lecture 5

Assignment on Traffic Networks

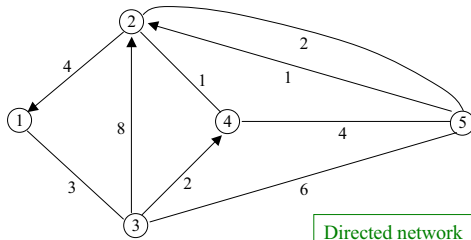
Prof. Ismail Chabini and Prof. Amedeo R. Odoni

Lecture 5 Outline

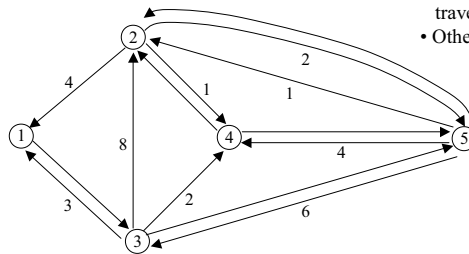
- Summary from previous lectures:
 - Assignment on non-congested networks: All-or-nothing assignment
 - Volume-delay functions for “congested” networks
- Framework for static traffic assignment models
- Static traffic assignment: concepts
- Static traffic assignment: principles
- User Optimal (UO) and System Optimal (SO) static traffic assignment
- Summary

Non-Congested Road Network and O-D Matrix

Mixed network



Directed network



O-D Matrix

	To	1	2	3	4	5
From 1	[-	30	35	40	15]
2	[10	-	15	12	10]
3	[50	40	-	35	20]
4	[25	30	35	-	40]
5	[45	30	35	40	-]

$[t(i,j)] = 3$

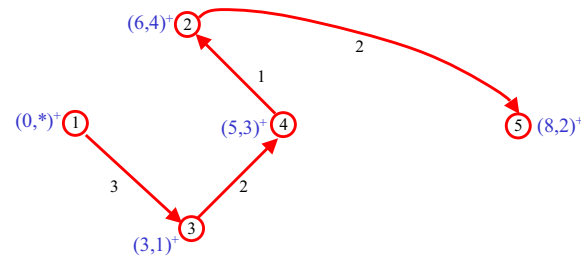
- In R7, $t(i,j)$ means demand for O-D pair (i,j) . ($t(i,j)$ does not mean travel time of link (i,j))
- Other common notation: $q(i,j)$ or q_{ij}

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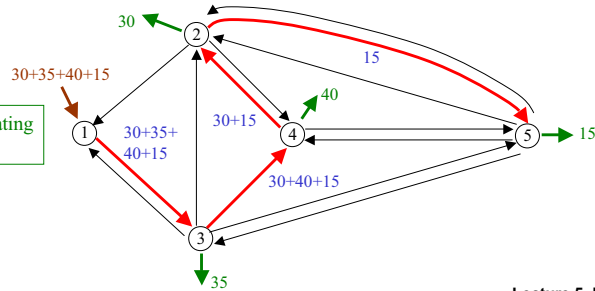
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Assign O-D Flows Originating from Node 1

Shortest path tree from node 1



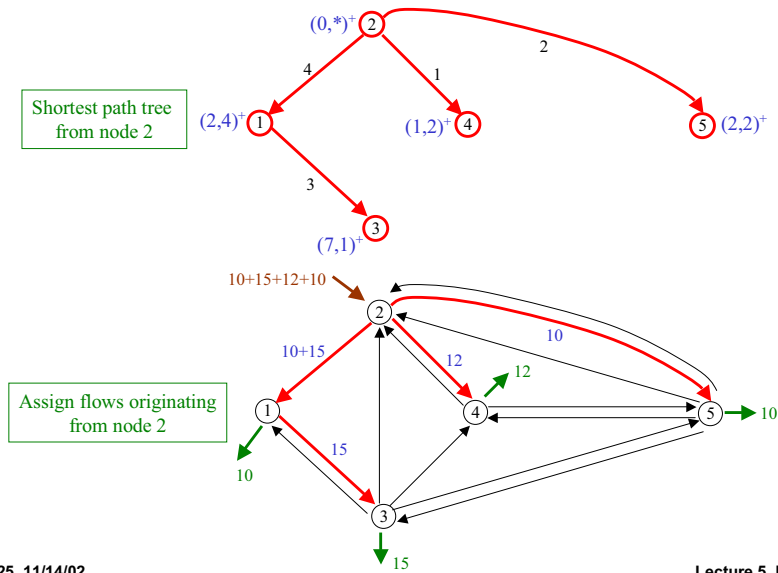
Assign flows originating from node 1



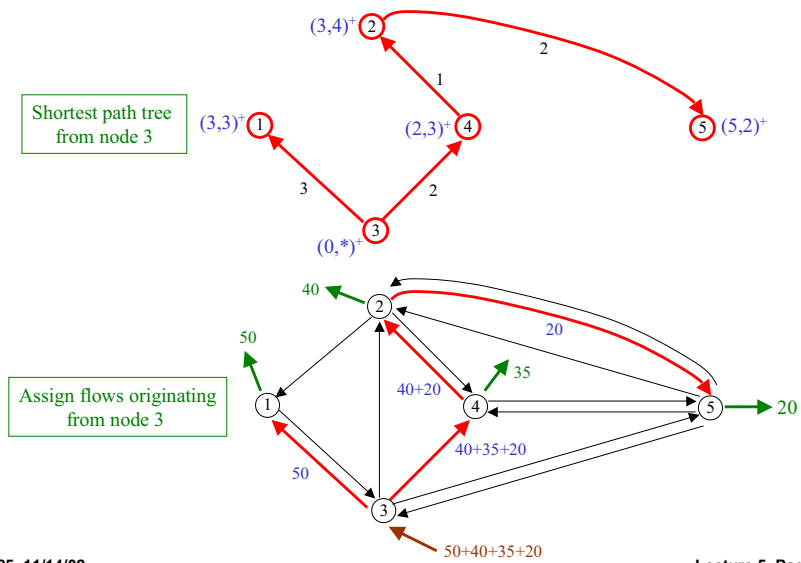
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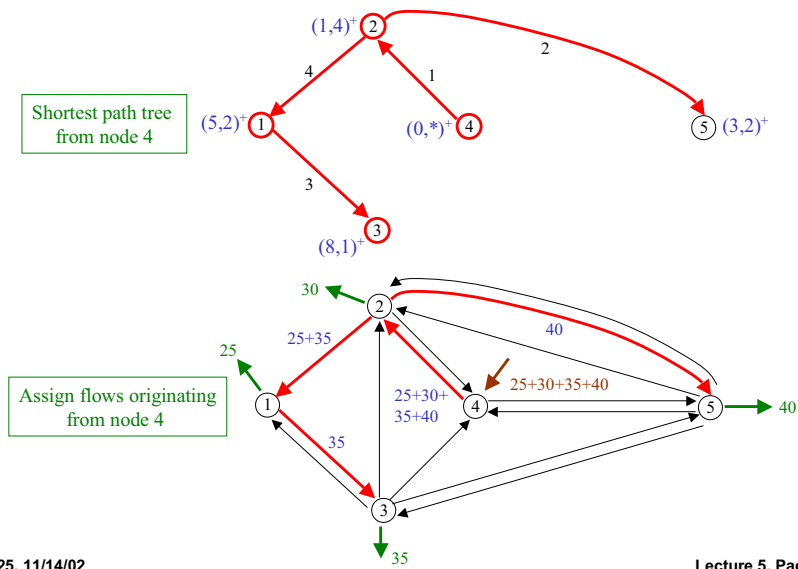
Assign O-D Flows Originating from Node 2



Assign O-D Flows Originating from Node 3



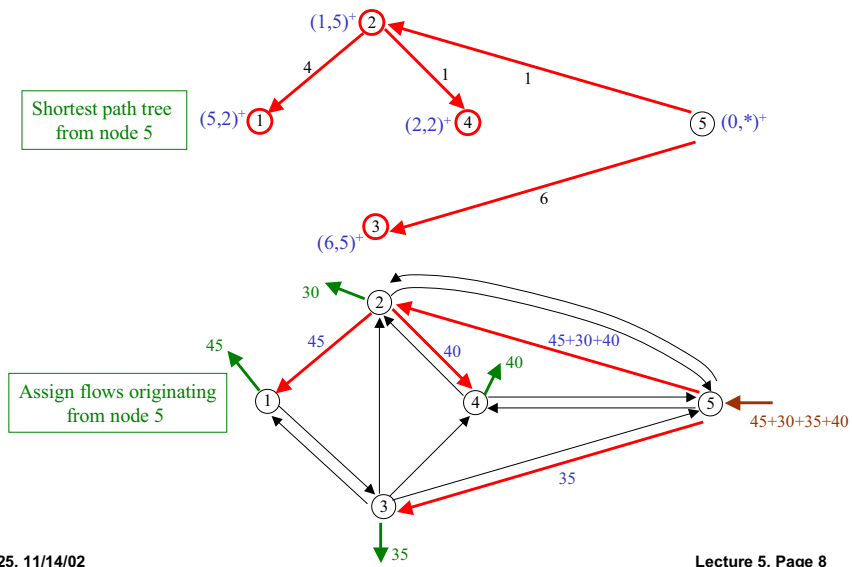
Assign O-D Flows Originating from Node 4



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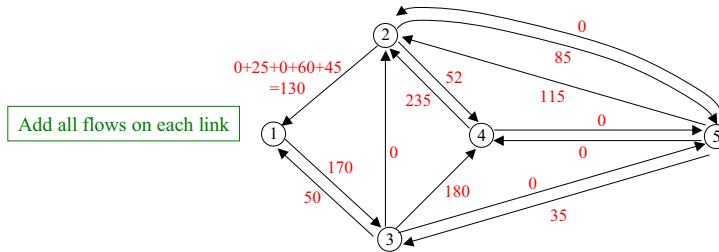
Assign O-D Flows Originating from Node 5



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Results of All-or-Nothing Assignment

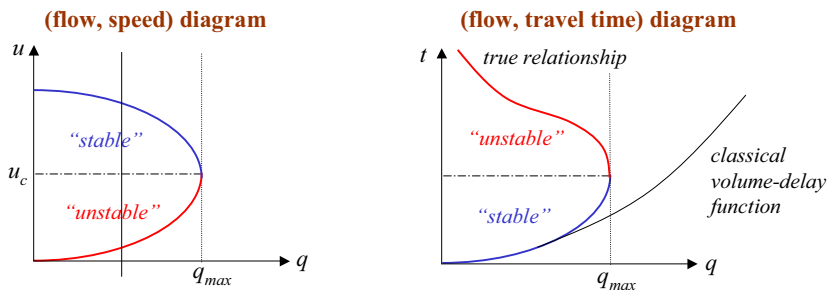


- All-or-nothing (AON) assignment** does not consider congestion
- The solution may not be unique (Why? Is unique solution important?)
- AON assignment does not make sense if certain links are congested
- To account for congestion, travel time must depend on link flow
- How to change the AON assignment method to make it work?
 - ⇒ **Equilibrium traffic assignment**

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Derived Diagrams from the Fundamental Diagram



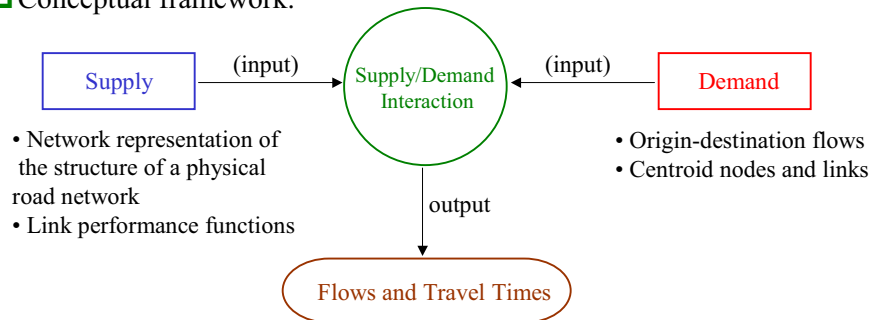
- In general, q cannot be used as a variable (why?)
- In the traffic planning area:
 - q is also called **volume**
 - travel time is also called **travel delay**
 - In the case of volume-delay functions, q is used as a variable

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Framework for Static Traffic Assignment Models

□ Conceptual framework:



- Network representation of the structure of a physical road network
- Link performance functions

- Origin-destination flows
- Centroid nodes and links

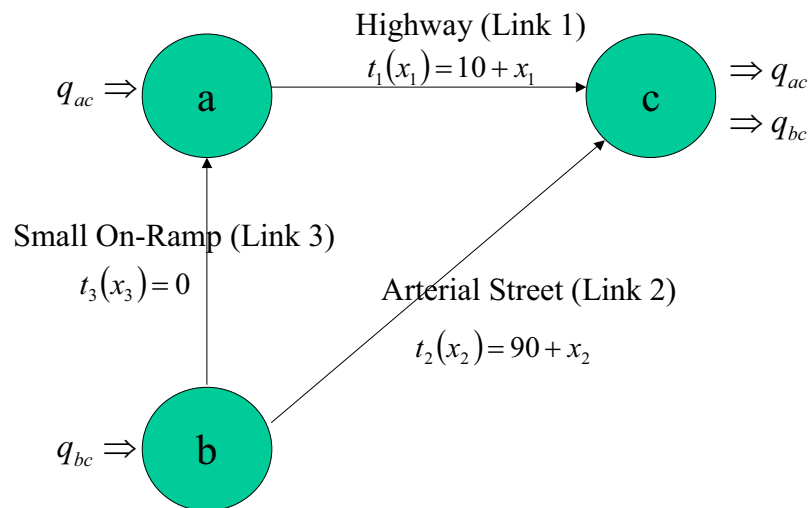
□ Principles of assignment to represent the supply/demand interaction

- *User Optimal (U.O.)*: O-D flows are assigned to paths with minimum travel time
- *System Optimal (S.O.)*: O-D flows are assigned such that total travel time on the network is minimum

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Example

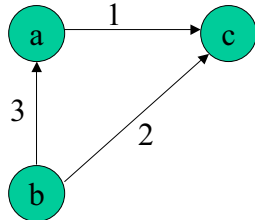


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Traffic Assignment Concepts

Conceptual Network



Demand

- O - Ds : (a, c) and (b, c)
- (a, c) : q_{ac} vehicles / hr
- (b, c) : q_{bc} vehicles / hr

Path flow variables

- O - D (a, c) : one path
1: Link 1 (f_1^{ac})
- O - D (b, c) : two paths
1: Link 2 (f_1^{bc})
2: Link 3, Link 1 (f_2^{bc})

O-D flows and path flows

$$q_{ac} = f_1^{ac}$$

$$q_{bc} = f_1^{bc} + f_2^{bc}$$

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Traffic Assignment Concepts (cont.)-1

Link (arc) flows and path flows

$$x_1 = f_1^{ac} + f_2^{bc}$$

$$x_2 = f_1^{bc}$$

$$x_3 = f_2^{bc}$$

Arc-path incidence matrix

O-D	→	a - c	b - c
Path	→	1	1 2
Link		$\begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$	

Assume that $f_2^{bc} = pq_{bc}$

$$x_1 = q_{ac} + pq_{bc}$$

$$x_2 = (1-p)q_{bc}$$

$$x_3 = pq_{bc}$$

Assignment matrix

$$\begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 1 & p \\ 0 & 1-p \\ 0 & p \end{bmatrix} \begin{bmatrix} q_{ac} \\ q_{bc} \end{bmatrix}$$

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Traffic Assignment Concepts (cont.)-2

□ t_1, t_2, t_3 are the travel times of links 1, 2, 3.

□ “Congested” networks:

- Link travel times depend on link flows
- Example:

$$t_1(x_1) = 10 + x_1, \quad t_2(x_2) = 90 + x_2, \quad t_3(x_3) = 0$$

□ Path travel-times as a function of link travel-times:

$$C_1^{ac} = t_1, \quad C_1^{bc} = t_2, \quad C_2^{bc} = t_1 + t_3$$

□ Total travel times (What is the unit of this quantity?):

$$x_1 * (10 + x_1) + x_2 * (90 + x_2) + x_3 * (0)$$

Assignment Principles

□ If OD flows are infinitely divisible:

- There is an infinite number of assignments
- Which assignment should we choose?
⇒ We need additional assumptions to define a less ambiguous assignment

□ Assignment principle: a principle used to determine an assignment

□ Examples of assignment principles:

- **User-optimal**: between each O-D pair, all used paths have equal and minimum travel times
- **System-optimal**: the total travel times are minimum

Mathematical Expressions of Assignment Principles

□ User-optimal traffic assignment principle: find p such that:

• O-D (b, c):

$$\text{If } p = 0, \quad (t_1 + t_3) \geq t_2$$

$$\text{If } p = 1, \quad (t_1 + t_3) \leq t_2$$

$$\text{If } 0 < p < 1, \quad t_1 + t_3 = t_2$$

• O-D (a, c): the question is not posed as there is only one path

□ System optimal: find p that minimizes:

$$f_1^{ac} C_1^{ac} + f_1^{bc} C_1^{bc} + f_2^{ac} C_2^{ac} (= x_1 t_1 + x_2 t_2 + x_3 t_3)$$

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Solution of the U.O. Assignment

□ We want to solve for a $p \in [0,1]$ such that:

$$\text{If } p = 0, \quad (t_1(x_1) + t_3(x_3)) \geq t_2(x_2)$$

$$\text{If } p = 1, \quad (t_1(x_1) + t_3(x_3)) \leq t_2(x_2)$$

$$\text{If } 0 < p < 1, \quad t_1(x_1) + t_3(x_3) = t_2(x_2)$$

and

$$\begin{array}{ll} q_{ac} (= 80) & t_1(x_1) = 10 + x_1 \\ q_{bc} (= 10) & t_2(x_2) = 90 + x_2 \\ & t_3(x_3) = 0 \end{array} \quad \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 1 & p \\ 0 & 1-p \\ 0 & p \end{bmatrix} \begin{bmatrix} q_{ac} \\ q_{bc} \end{bmatrix}$$

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Solution of the U.O. Assignment (cont.)

- $t_2(x_2) - (t_1(x_1) + t_3(x_3)) = (90 + x_2) - (10 + x_1) = 80 + x_2 - x_1$
- $x_2 - x_1 = (q_{bc} - x_3) - (q_{ac} + x_3) = (q_{bc} - q_{ac}) - 2x_3$
- $t_2(x_2) - (t_1(x_1) + t_3(x_3)) = 80 + q_{bc} - q_{ac} - 2x_3$
- If $\frac{80 + q_{bc} - q_{ac}}{2} \geq 0$, then $x_3 = \frac{80 + q_{bc} - q_{ac}}{2}$
- Example: $(q_{ac}, q_{bc}) = (80, 10) \rightarrow (x_1, x_2, x_3) = (85, 5, 5)$
- The travel time on any path is: 95
- Total travel time (per hour): $(80+10)*95=8550$ veh-hrs (per hour)

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(UO) Assignment May Depend on the Demand

- If $\frac{80 + q_{bc} - q_{ac}}{2} \geq 0$, then

$$(x_1^*, x_2^*, x_3^*) = \left(\frac{80 + q_{bc} + q_{ac}}{2}, \frac{q_{bc} - 80 + q_{ac}}{2}, \frac{80 + q_{bc} - q_{ac}}{2} \right)$$
- If $\frac{80 + q_{bc} - q_{ac}}{2} < 0$, then $(x_1^*, x_2^*, x_3^*) = (q_{ac}, q_{bc}, 0)$
- Proof: $p^* = 0$ and $t_2(q_{bc}) - (t_1(q_{ac}) - t_3(0)) = \frac{80 + q_{bc} - q_{ac}}{2} < 0$
- The U.O. assignment is a function of the demand, and the function may be non-linear. Example:

$$(q_{ac}, q_{bc}) = (80, 10) \rightarrow (x_1^*, x_2^*, x_3^*) = (85, 5, 5)$$

$$(q_{ac}, q_{bc}) = (160, 20) \rightarrow (x_1^*, x_2^*, x_3^*) = (160, 20, 0)$$

$$(160, 20) = 2 \times (80, 10) \text{ But } (160, 20, 0) \neq 2 \times (85, 5, 5)$$
- Remark: An increase in demand may lead to a decrease in the flow on a link

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Building More Roads Is Not Always Better

- Without Link 3, there is only one possible assignment
 $(x_1^*, x_2^*) = (q_{ac}, q_{bc}) = (80, 10)$
 Total travel times in one hour: $80 \cdot (10 + 80) + 10 \cdot (90 + 10) = 8200$ veh-hr
 (in one hour)
- U.O. with Link 3:
 Total travel times (in one hour): 8550 veh-hr (in one hour)
- System travel times are worse if one adds Link 3! (Is this intuitive?)
- This phenomenon is known as Braess “Paradox”, and is not an isolated phenomenon.

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System Optimal Assignment

- $$\min x_1 t_1(x_1) + x_2 t_2(x_2) + x_3 t_3(x_3)$$

$$s.t. \quad f_1^{ac} = q_{ac} \quad \text{Demand}$$

$$f_1^{bc} + f_2^{bc} = q_{bc}$$

$$f_1^{ac} \geq 0, f_1^{bc} \geq 0, f_2^{bc} \geq 0 \quad \text{Non-negativity}$$

$$x_1 = f_1^{ac} + f_1^{bc}$$

$$x_2 = f_1^{bc}$$

$$x_3 = f_2^{bc} \quad \text{Definition of link flows}$$
- S.O. solution: $(x_1^*, x_2^*, x_3^*) = (80, 10, 0)$, if $(q_{ac}, q_{bc}) = (80, 10)$
- $$\min x_1 t_1(x_1) + x_2 t_2(x_2) + x_3 t_3(x_3) = \min \int_0^{x_1} m_1(x_1) dx_1 + \int_0^{x_2} m_2(x_2) dx_2 + \int_0^{x_3} m_3(x_3) dx_3$$

$$\text{Where: } m_1(x_1) = \frac{d(x_1 t_1(x_1))}{dx_1}, m_2(x_2) = \frac{d(x_2 t_2(x_2))}{dx_2}, m_3(x_3) = \frac{d(x_3 t_3(x_3))}{dx_3}$$

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Lecture 5 Summary

- Assignment on non-congested networks: All-or-nothing assignment
- Volume-delay functions for “congested” networks and static demand
- Framework for static traffic assignment models
- Static traffic assignment concepts and principles
- User Optimal and System Optimal static traffic assignment
- Summary