## Constraint Programming: Modeling, Arc Consistency and Propagation

Brian C. Williams
16.410-13

September 22nd, 2010
Slides draw material from:
6.034 notes, by Tomas Lozano Perez

AIMA, by Stuart Russell \& Peter Norvig
Constraint Processing, by Rina Dechter

## Assignments

- Assignment:
- Problem Set \#2 due today, Wed. Sept. 22 ${ }^{\text {nd }}, 2010$.
- Problem Set \#3: Analysis, Path Planning and Constraint Programming, out today, due Wed., Sept. 29 ${ }^{\text {th }}, 2010$.
- Reading:
- Today: [AIMA] Ch. 6.1, 24.3-5; Constraint Modeling.
- Monday: [AIMA] Ch. 6.2-5; Constraint Satisfaction.
- To Learn More: Constraint Processing, by Rina Dechter
- Ch. 2: Constraint Networks
-Ch. 3: Consistency Enforcing and Propagation


## Outline

- Interpreting line diagrams
- Constraint satisfaction problems (CSP) [aka constraint programs (CP)].
- Solving CSPs
- Case study: Scheduling (Appendix)


## Outline

- Interpreting line diagrams
- Constraint modeling
- Constraint propagation
- Constraint satisfaction problems (CSP) aka constraint programs (CP)
- Solving CSPs
- Case study: Scheduling (Appendix)



## Line Labeling as Constraint Programming



18 vertex labelings that are


Huffman Clowes (1971):
Interpretation of opaque, trihedral solids with no surface marks.
Waltz (1972): Compute labeling through local propagation. ,

## Outline

- Interpreting line diagrams
- Constraint modeling
- Constraint propagation
- Constraint satisfaction problems (CSP) aka constraint programs (CP).
- Solving CSPs
- Case study: Scheduling (Appendix)


## Modeling: Make Simplifying Assumptions

1. Limited line interpretations:

No shadows or cracks.
2. Three-faced vertices:

Intersection of exactly three object faces
(e.g., no pyramid tops).
3. General position:

Small perturbations of selected viewing points can not lead to a change in junction type.

Consider:

## Modeling: Systematically derive all realizable junction types

- a three face vertex, which divides space into octants,
- (not guaranteed to be at right angles), and
- all possible fillings of octants, viewed from all empty octants.



## Modeling: Systematically derive all realizable junction types

- Case 1: View seven filled octants from the only empty octant.



## Modeling: Systematically derive all realizable junction types

- Case 2a: View one filled octant from all empty upper octants....


Brian Williams, Fall 10

## Modeling: Systematically derive all realizable junction types

- Case 2b: View one filled octant from all empty lower octants.



## Outline

- Interpreting line diagrams
- Constraint modeling
- Constraint propagation
- Constraint satisfaction problems (CSP) aka constraint programs (CP).
- Solving CSPs
- Case study: Scheduling (Appendix)



Without background borders, interpretations become unstable.


Brian Williams, Fall 10

## Outline

- Interpreting line diagrams
- Constraint satisfaction problems (CSP) aka constraint programs (CP).
- Solving CSPs
- Case study: Scheduling (appendix)


## Constraint Satisfaction Problems

4 Queens Problem:
Place 4 queens on a $4 \times 4$ chessboard so that no queen can attack another.


How do we formulate?
Variables Chessboard positions

Domains Queen 1-4 or blank
Constraints Two positions on a line (vertical, horizontal, diagonal) cannot both be Q


## Constraint Satisfaction Problems (CSP)

Input: A Constraint Satisfaction Problem is a triple $<\mathrm{V}, \mathrm{D}, \mathrm{C}>$, where:

- $V$ is a set of variables $V_{i}$
- $D$ is a set of variable domains,
- The domain of variable $V_{i}$ is denoted $D_{i}$
- $C=$ is a set of constraints on assignments to $V$
- Each constraint $\mathrm{C}_{\mathrm{i}}=\left\langle\mathrm{S}_{\mathrm{i}}, \mathrm{R}_{\mathrm{i}}\right\rangle$ specifies allowed variable assignments.
- $S_{i}$ the constraint's scope, is a subset of variables V .
- $R_{i}$ the constraint's relation, is a set of assignments to $S_{i}$.

Output: A full assignment to V , from elements of V 's domain, such that all constraints in $C$ are satisfied.

Example: "Provide one A and two B's."

- $V=\{A, B\}$, each with domain $D_{i}=\{1,2\}$
- $C=\{<\{A, B\},\{<1,2>,<1,1>\}>$
$<\{A, B\},\{<1,2>,<2,2>\}>\}$
- Output: <1,2>

Brian Williams, Fall 10
"one A"
"two Bs"
(for example)

## Conventions

- List scope in subscript.
- Specify one constraint per scope.

Example: "Provide one A and two B's."

- $C=\left\{C_{A B}\right\}$

$$
C_{A B}=\{<1,2>\}
$$

- $C=\left\{C_{A}, C_{B}\right\}$

$$
C_{A}=\{<1>\}
$$

$$
C_{B}=\{<2>\}
$$

## Good Encodings Are Essential: 4 Queens

4 Queens Problem:
Place 4 queens on a $4 \times 4$
chessboard so that no queen can attack another.


How big is the encoding?
Variables Chessboard positions

Domains Queen 1-4 or blank
Constraints Two positions on a line (vertical, horizontal, diagonal) cannot both be Q

## Good Encodings Are Essential: 4 Queens

Place queens so that no queen can attack another.

What is a better encoding?

|  |  |  | $Q$ |  |
| :--- | :--- | :--- | :--- | :--- |
|  |  |  |  |  |
|  |  |  |  | $Q$ |
|  | $Q$ |  |  |  |
|  |  |  |  | $Q$ |

- Assume one queen per column.
- Determine what row each queen should be in.

Variables $\quad Q_{1}, Q_{2}, Q_{3}, Q_{4}$,
Domains $\quad\{1,2,3,4\}$
Constraints $Q_{i}<>Q_{j} \quad$ "On different rows" $\left|Q_{i}-Q_{j}\right|<>|i-j| \quad$ "Stay off the diagonals"

Example

$$
C_{1,2}=\{(1,3)(1,4)(2,4)(3,1)(4,1)(4,2)\}
$$

## Good Encodings Are Essential: 4 Queens

Place queens so that no queen can attack another.

Variables $\quad Q_{1}, Q_{2}, Q_{3}, Q_{4}$,

| 1 | $\uparrow$ | $Q$ |  |  |
| :--- | :--- | :--- | :--- | :--- |
|  |  |  |  |  |
|  |  |  |  | $Q$ |
|  | $Q$ |  |  |  |
|  |  |  |  |  |
|  |  |  | $Q$ |  |

Domains $\quad\{1,2,3,4\}$
Constraints $Q_{i}<>Q_{j} \quad$ "On different rows"

$$
\left|Q_{i}-Q_{j}\right|<>|i-j| \quad \text { "Stay off the diagonals" }
$$

Example: $\quad \mathrm{C}_{1,2}=\{(1,3)(1,4)(2,4)(3,1)(4,1)(4,2)\}$
What is $\mathrm{C}_{13}$ ?

## A general class of CSPs

Finite Domain, Binary CSPs

- each constraint relates at most two variables.
- each variable domain is finite.

Property: all n-ary CSPs are reducible to binary CSPs.

Depict as a Constraint Graph

- Nodes (vertices) are variables.
- Arcs (edges) are binary constraints.



## Example: Graph Coloring

Pick colors for map regions, without coloring adjacent regions with the same color

Variables regions

Domains
allowed colors


Constraints adjacent regions must have different colors

## Outline

- Interpreting line problems
- Constraint satisfaction problems (CSP) aka constraint programs (CP).
- Solving CSPs
- Arc-consistency and propagation
- Analysis of constraint propagation (next lecture)
- Search (next lecture)
- Case study: Scheduling (appendix)


## Good News / Bad News

Good News

- very general \& interesting family of problems.
- Problem formulation used extensively in autonomy and decision making applications.

Bad News includes NP-Hard (intractable ?) problems

## Algorithmic Design Paradigm

Solving CSPs involves a combination of:

1. Inference

- Solve partially by eliminating values that can't be part of any solution (constraint propagation).
- Make implicit constraints explicit.

2. Search

- Try alternative assignments against constraints.

Inference: Waltz constraint propagation for visual interpretation generalizes to arc-consistency and the AC-3 algorithm.

## Directed Arc Consistency

Idea: Eliminate values of a variable domain that can never satisfy a specified constraint (an arc).


Definition: $\operatorname{arc}\left\langle\mathrm{x}_{\mathrm{i}}, \mathrm{x}_{\mathrm{j}}\right\rangle$ is arc consistent if $\left\langle\mathrm{x}_{\mathrm{i}}, \mathrm{x}_{\mathrm{j}}\right\rangle$ and $\left\langle\mathrm{x}_{\mathrm{j}}, \mathrm{x}_{\mathrm{i}}\right\rangle$ are directed arc consistent.


## Directed Arc Consistency



Definition: arc $<x_{i}, x_{j}>$ is directed arc consistent if

- for every $a_{i}$ in $D_{i}$,
- there exists some $a_{j}$ in $D_{j}$ such that
- assignment $<\mathrm{a}_{\mathrm{i}}, \mathrm{a}_{\mathrm{j}}>$ satisfies constraint $\mathrm{C}_{\mathrm{ij}}$,
$\cdot \forall a_{i} \in D_{i}, \exists a_{j} \in D_{j}$ such that $\left.<a_{i}, a_{j}\right\rangle \in C_{i j}$
- $\forall$ denotes "for all," $\exists$ denotes "there exists" and $\in$ denotes "in." Brian Williams, Fall 10


## Revise: A directed arc consistency procedure

Definition: arc $\left\langle x_{i}, x_{j}\right\rangle$ is directed arc consistent if $\forall \mathrm{a}_{\mathrm{i}} \in \mathrm{D}_{\mathrm{i}}, \exists \mathrm{a}_{\mathrm{j}} \in \mathrm{D}_{\mathrm{j}}$ such that $\left\langle\mathrm{a}_{\mathrm{i}}, \mathrm{a}_{\mathrm{j}}\right\rangle \in \mathrm{C}_{\mathrm{ij}}$

Revise ( $\mathrm{x}_{\mathrm{i}}, \mathrm{x}_{\mathrm{j}}$ )
Input: Variables $x_{i}$ and $x_{j}$ with domains $D_{i}$ and $D_{j}$ and constraint relation $R_{i j}$. Output: pruned $D_{i}$, such that $x_{i}$ is directed arc-consistent relative to $x_{j}$.

1. for each $a_{i} \in D_{i}$
2. if there is no $a_{j} \in D_{j}$ such that $\left\langle a_{i}, a_{j}\right\rangle \in R_{i j}$
3. then delete $\mathrm{a}_{\mathrm{i}}$ from $\mathrm{D}_{\mathrm{i}}$.
4. endif
5. endfor

Constraint Processing, by R. Dechter pgs 54-6

## Full Arc Consistency over All Constraints via Constraint Propagation

Definition: arc $\left.<x_{i}, x_{j}\right\rangle$ is directed arc consistent if $\forall a_{i} \in D_{i}, \exists a_{j} \in D_{j}$ such that $\left\langle a_{i}, a_{j}\right\rangle \in C_{i j}$

Constraint Propagation:
To achieve (directed) arc consistency over CSP:

1. For every arc $C_{i j}$ in CSP, with tail domain $D_{i}$, call Revise.
2. Repeat until quiescence:

If an element was deleted from $D_{i}$, then repeat Step 1
(AC-1)

## Full Arc-Consistency via AC-1

## AC-1(CSP)

Input: A constraint satisfaction problem CSP = <X, D, C>. Output: CSP', the largest arc-consistent subset of CSP.

1. repeat
2. for every $\mathrm{c}_{\mathrm{ij}} \in \mathrm{C}$,
3. Revise $\left(\mathrm{x}_{\mathrm{i}}, \mathrm{x}_{\mathrm{i}}\right)$
4. Revise $\left(\mathrm{x}_{\mathrm{j}}, \mathrm{x}_{\mathrm{i}}\right)$
5. endfor

For every arc, prune head and tail domains.

Constraint Processing,
6. until no domain is changed.
by R. Dechter
pgs 57

## Full Arc Consistency via Constraint Propagation

Definition: arc $<x_{i}, x_{j}>$ is directed arc consistent if

$$
\forall a_{i} \in D_{i}, \exists a_{j} \in D_{j} \text { such that }<a_{i}, a_{j}>\in C_{i j}
$$

Constraint Propagation:
To achieve (directed) arc consistency over CSP:

1. For every arc $\mathrm{C}_{\mathrm{ij}}$ in CSP, with tail domain $\mathrm{D}_{\mathrm{i}}$, call Revise.
2. Repeat until quiescence:

If an element was deleted from $D_{i}$, then

$$
\text { repeat Step } 1
$$

OR call Revise on each arc with head $D_{i}$
(use FIFO Q, remove duplicates)

## Full Arc-Consistency via AC-3 (Waltz CP)

## AC-3(CSP)

Input: A constraint satisfaction problem CSP = <X, D, C>. Output: CSP', the largest arc-consistent subset of CSP.

1. for every $\mathrm{c}_{\mathrm{ij}} \in \mathrm{C}$,

Constraint Processing,
2. queue $\leftarrow$ queue $\cup\left\{\left\langle\mathrm{x}_{\mathrm{i}}, \mathrm{x}_{\mathrm{j}}\right\rangle,\left\langle\mathrm{x}_{\mathrm{j}}, \mathrm{x}_{\mathrm{i}}\right\rangle\right\}$
3. endfor pgs 58-9
4. while queue $\neq\{ \}$
5. select and delete arc $\left\langle x_{i}, x_{j}\right\rangle$ from queue
6. $\operatorname{Revise}\left(\mathrm{x}_{\mathrm{i}}, \mathrm{x}_{\mathrm{j}}\right)$
7. if Revise $\left(x_{i}, x_{j}\right)$ caused a change in $D_{i}$
8. then queue $\leftarrow$ queue $\cup\left\{\left\langle x_{k}, x_{i}\right\rangle \mid k \neq i, k \neq j\right\}$
9. endif
10. endwhile

## Constraint Propagation Example AC-3

## Graph Coloring

Initial Domains


Each undirected arc denotes two directed arcs.

## Constraint Propagation Example AC-3

Graph Coloring
Initial Domains


| Arc examined | Value deleted |
| :--- | :--- |
|  |  |
|  |  |
|  |  |
|  |  |
|  |  |



Arcs to examine
$\mathrm{V}_{1}-\mathrm{V}_{2}, \mathrm{~V}_{1}-\mathrm{V}_{3}, \mathrm{~V}_{2}-\mathrm{V}_{3}$

- Introduce queue of arcs to be examined.
- Start by adding all arcs to the queue.


## Constraint Propagation Example AC-3



| Arc examined | Value deleted |
| :--- | :--- |
|  |  |
|  |  |
|  |  |
|  |  |
|  |  |
|  |  |



Arcs to examine

$$
V_{1}-V_{2}, V_{1}-V_{3}, V_{2}-V_{3}
$$

- $\mathrm{V}_{\mathrm{i}}-\mathrm{V}_{\mathrm{j}}$ denotes two arcs, between $\mathrm{V}_{\mathrm{i}}$ and $\mathrm{V}_{\mathrm{j}}$.
- $\mathrm{V}_{\mathrm{i}}>\mathrm{V}_{\mathrm{j}}$ denotes an arc from $\mathrm{V}_{\mathrm{i}}$ to $\mathrm{V}_{\mathrm{j}}$. 42


## Constraint Propagation Example AC-3

Graph Coloring
Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $\mathrm{V}_{1}>\mathrm{V}_{2}$ |  |
|  |  |
|  |  |
|  |  |
|  |  |
|  |  |



Arcs to examine

$$
V_{2}>V_{1}, V_{1}-V_{3}, V_{2}-V_{3}
$$

- Delete unmentioned tail values $\cdot \mathrm{V}_{\mathrm{i}}-\mathrm{V}_{\mathrm{j}}$ denotes two arcs, between $\mathrm{V}_{\mathrm{i}}$ and $\mathrm{V}_{\mathrm{j}}$.
- $\mathrm{V}_{\mathrm{i}}>\mathrm{V}_{\mathrm{j}}$ denotes an arc from $\mathrm{V}_{\mathrm{i}}$ to $\mathrm{V}_{\mathrm{j}}$. 43


## Constraint Propagation Example AC-3

## Graph Coloring

Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $\mathrm{V}_{1}>\mathrm{V}_{2}$ | none |
|  |  |
|  |  |
|  |  |
|  |  |
|  |  |



Arcs to examine

$$
\mathrm{V}_{2}>\mathrm{V}_{1}, \mathrm{~V}_{1}-\mathrm{V}_{3}, \mathrm{~V}_{2}-\mathrm{V}_{3}
$$

- Delete unmentioned tail values $\cdot \mathrm{V}_{\mathrm{i}}-\mathrm{V}_{\mathrm{j}}$ denotes two arcs, between $\mathrm{V}_{\mathrm{i}}$ and $\mathrm{V}_{\mathrm{j}}$.
- $\mathrm{V}_{\mathrm{i}}>\mathrm{V}_{\mathrm{j}}$ denotes an arc from $\mathrm{V}_{\mathrm{i}}$ to $\mathrm{V}_{\mathrm{j}}$. ${ }_{44}$


## Constraint Propagation Example AC-3

Graph Coloring
Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $\mathrm{V}_{1}>\mathrm{V}_{2}$ | none |
| $\mathrm{V}_{2}>\mathrm{V}_{1}$ |  |
|  |  |
|  |  |
|  |  |
|  |  |



Arcs to examine
$V_{1}-V_{3}, V_{2}-V_{3}$

- Delete unmentioned tail values $\cdot \mathrm{V}_{\mathrm{i}}-\mathrm{V}_{\mathrm{j}}$ denotes two arcs, between $\mathrm{V}_{\mathrm{i}}$ and $\mathrm{V}_{\mathrm{j}}$.
- $\mathrm{V}_{\mathrm{i}}>\mathrm{V}_{\mathrm{j}}$ denotes an arc from $\mathrm{V}_{\mathrm{i}}$ to $\mathrm{V}_{\mathrm{j}}$. ${ }^{45}$


## Constraint Propagation Example AC-3

## Graph Coloring

Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $\mathrm{V}_{1}>\mathrm{V}_{2}$ | none |
| $\mathrm{V}_{2}>\mathrm{V}_{1}$ | none |
|  |  |
|  |  |
|  |  |
|  |  |



Arcs to examine

$$
V_{1}-V_{3}, V_{2}-V_{3}
$$

- Delete unmentioned tail values $\cdot \mathrm{V}_{\mathrm{i}}-\mathrm{V}_{\mathrm{j}}$ denotes two arcs, between $\mathrm{V}_{\mathrm{i}}$ and $\mathrm{V}_{\mathrm{j}}$.
- $\mathrm{V}_{\mathrm{i}}>\mathrm{V}_{\mathrm{j}}$ denotes an arc from $\mathrm{V}_{\mathrm{i}}$ to $\mathrm{V}_{\mathrm{j}}$. ${ }_{\text {46 }}$


## Constraint Propagation Example AC-3

Graph Coloring
Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $\mathrm{V}_{1}-\mathrm{V}_{2}$ | none |
|  |  |
|  |  |
|  |  |
|  |  |
|  |  |



Arcs to examine
$\mathrm{V}_{1}-\mathrm{V}_{3}, \mathrm{~V}_{2}-\mathrm{V}_{3}$

- Delete unmentioned tail values $\cdot \mathrm{V}_{\mathrm{i}}-\mathrm{V}_{\mathrm{j}}$ denotes two arcs, between $\mathrm{V}_{\mathrm{i}}$ and $\mathrm{V}_{\mathrm{j}}$.
- $\mathrm{V}_{\mathrm{i}}>\mathrm{V}_{\mathrm{j}}$ denotes an arc from $\mathrm{V}_{\mathrm{i}}$ to $\mathrm{V}_{\mathrm{j}}$. ${ }_{47}$


## Constraint Propagation Example AC-3

## Graph Coloring

Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $\mathrm{V}_{1}-\mathrm{V}_{2}$ | none |
| $\mathrm{V}_{1}>\mathrm{V}_{3}$ |  |
|  |  |
|  |  |
|  |  |
|  |  |



Arcs to examine

$$
V_{3}>V_{1}, V_{2}-V_{3}
$$

- Delete unmentioned tail values $\cdot \mathrm{V}_{\mathrm{i}}-\mathrm{V}_{\mathrm{j}}$ denotes two arcs, between $\mathrm{V}_{\mathrm{i}}$ and $\mathrm{V}_{\mathrm{j}}$.
- $\mathrm{V}_{\mathrm{i}}>\mathrm{V}_{\mathrm{j}}$ denotes an arc from $\mathrm{V}_{\mathrm{i}}$ to $\mathrm{V}_{\mathrm{j}}$. ${ }^{48}$


## Constraint Propagation Example AC-3

Graph Coloring
Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $V_{1}-V_{2}$ | none |
| $V_{1}>V_{3}$ | $V_{1}(G)$ |
|  |  |
|  |  |
|  |  |
|  |  |



Arcs to examine

$$
V_{3}>V_{1}, V_{2}-V_{3}
$$

IF An element of a variable's domain is removed,

## Constraint Propagation Example AC-3

## Graph Coloring

Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $V_{1}-V_{2}$ | none |
| $V_{1}>V_{3}$ | $V_{1}(\mathbf{G})$ |
|  |  |
|  |  |
|  |  |
|  |  |



IF An element of a variable's domain is removed, THEN add all arcs to that variable to the examination queue.

## Constraint Propagation Example AC-3

Graph Coloring
Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $V_{1}-V_{2}$ | none |
| $V_{1}>V_{3}$ | $V_{1}(G)$ |
| $V_{3}>V_{1}$ |  |
|  |  |
|  |  |
|  |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed,
THEN add all arcs to that variable to the examination queue. 51

## Constraint Propagation Example AC-3

## Graph Coloring

Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $V_{1}-V_{2}$ | none |
| $V_{1}>V_{3}$ | $V_{1}(G)$ |
| $V_{3}>V_{1}$ | none |
|  |  |
|  |  |
|  |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed, THEN add all arcs to that variable to the examination queue. 52

## Constraint Propagation Example AC-3

Graph Coloring
Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $\mathrm{V}_{1}-\mathrm{V}_{2}$ | none |
| $\mathrm{V}_{1}-\mathrm{V}_{3}$ | $\mathrm{~V}_{1}(\mathrm{G})$ |
|  |  |
|  |  |
|  |  |
|  |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed,
THEN add all arcs to that variable to the examination queue.

## Constraint Propagation Example AC-3

## Graph Coloring

Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $\mathrm{V}_{1}-\mathrm{V}_{2}$ | none |
| $\mathrm{V}_{1}-\mathrm{V}_{3}$ | $\mathrm{~V}_{1}(\mathrm{G})$ |
| $\mathrm{V}_{2}>\mathrm{V}_{3}$ |  |
|  |  |
|  |  |
|  |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed, THEN add all arcs to that variable to the examination queue. 54

## Constraint Propagation Example AC-3

## Graph Coloring

Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $\mathrm{V}_{1}-\mathrm{V}_{2}$ | none |
| $\mathrm{V}_{1}-\mathrm{V}_{3}$ | $\mathrm{~V}_{1}(\mathrm{G})$ |
| $\mathrm{V}_{2}>\mathrm{V}_{3}$ | $\mathrm{~V}_{2}(G)$ |
|  |  |
|  |  |
|  |  |



Arcs to examine

- Delete unmentioned tail values

IF An element of a variable's domain is removed,
THEN add all arcs to that variable to the examination queue.

## Constraint Propagation Example AC-3

## Graph Coloring

Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $\mathrm{V}_{1}-\mathrm{V}_{2}$ | none |
| $\mathrm{V}_{1}-\mathrm{V}_{3}$ | $\mathrm{~V}_{1}(\mathrm{G})$ |
| $\mathrm{V}_{2}>\mathrm{V}_{3}$ | $\mathrm{~V}_{2}(\mathrm{G})$ |
|  |  |
|  |  |
|  |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed, THEN add all arcs to that variable to the examination queue. ${ }_{56}^{56}$

## Constraint Propagation Example AC-3

Graph Coloring
Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $V_{1}-V_{2}$ | none |
| $V_{1}-V_{3}$ | $V_{1}(G)$ |
| $V_{2}>V_{3}$ | $V_{2}(G)$ |
|  |  |
|  |  |
|  |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed,
THEN add all arcs to that variable to the examination queue. 57

## Constraint Propagation Example AC-3

## Graph Coloring

Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $V_{1}-V_{2}$ | none |
| $V_{1}-V_{3}$ | $V_{1}(G)$ |
| $V_{2}>V_{3}$ | $V_{2}(G)$ |
| $V_{3}>V_{2}$ |  |
|  |  |
|  |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed, THEN add all arcs to that variable to the examination queue. 58

## Constraint Propagation Example AC-3

Graph Coloring
Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $V_{1}-V_{2}$ | none |
| $V_{1}-V_{3}$ | $V_{1}(G)$ |
| $V_{2}>V_{3}$ | $V_{2}(G)$ |
| $V_{3}>V_{2}$ | none |
|  |  |
|  |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed,
THEN add all arcs to that variable to the examination queue.

## Constraint Propagation Example AC-3

## Graph Coloring

Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $\mathrm{V}_{1}-\mathrm{V}_{2}$ | none |
| $\mathrm{V}_{1}-\mathrm{V}_{3}$ | $\mathrm{~V}_{1}(\mathrm{G})$ |
| $\mathrm{V}_{2}-\mathrm{V}_{3}$ | $\mathrm{~V}_{2}(\mathrm{G})$ |
|  |  |
|  |  |
|  |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed, THEN add all arcs to that variable to the examination queue. 60

## Constraint Propagation Example AC-3

Graph Coloring
Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $V_{1}-V_{2}$ | none |
| $V_{1}-V_{3}$ | $V_{1}(G)$ |
| $V_{2}-V_{3}$ | $V_{2}(G)$ |
| $V_{2}>V_{1}$ |  |
|  |  |
|  |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed,
THEN add all arcs to that variable to the examination queue.
61

## Constraint Propagation Example AC-3

## Graph Coloring

Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $\mathrm{V}_{1}-\mathrm{V}_{2}$ | none |
| $\mathrm{V}_{1}-\mathrm{V}_{3}$ | $\mathrm{~V}_{1}(\mathrm{G})$ |
| $\mathrm{V}_{2}-\mathrm{V}_{3}$ | $\mathrm{~V}_{2}(\mathrm{G})$ |
| $\mathrm{V}_{2}>\mathrm{V}_{1}$ | none |
|  |  |
|  |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed, THEN add all arcs to that variable to the examination queue. 62

## Constraint Propagation Example AC-3

Graph Coloring
Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $V_{1}-V_{2}$ | none |
| $V_{1}-V_{3}$ | $V_{1}(G)$ |
| $V_{2}-V_{3}$ | $V_{2}(G)$ |
| $V_{2}>V_{1}$ | none |
| $V_{1}>V_{2}$ |  |
|  |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed,
THEN add all arcs to that variable to the examination queue.

## Constraint Propagation Example AC-3



| Arc examined | Value deleted |
| :---: | :---: |
| $V_{1}-V_{2}$ | none |
| $V_{1}-V_{3}$ | $V_{1}(G)$ |
| $V_{2}-V_{3}$ | $V_{2}(G)$ |
| $V_{2}>V_{1}$ | none |
| $V_{1}>V_{2}$ | $V_{1}(R)$ |
|  |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed, THEN add all arcs to that variable to the examination queue. ${ }_{64}$

## Constraint Propagation Example AC-3

Graph Coloring
Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $V_{1}-V_{2}$ | none |
| $V_{1}-V_{3}$ | $V_{1}(G)$ |
| $V_{2}-V_{3}$ | $V_{2}(G)$ |
| $V_{2}>V_{1}$ | none |
| $V_{1}>V_{2}$ | $V_{1}(R)$ |
|  |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed,
THEN add all arcs to that variable to the examination queue.
65

## Constraint Propagation Example AC-3

## Graph Coloring

Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $V_{1}-V_{2}$ | none |
| $V_{1}-V_{3}$ | $V_{1}(G)$ |
| $V_{2}-V_{3}$ | $V_{2}(G)$ |
| $V_{2}-V_{1}$ | $V_{1}(R)$ |
|  |  |
|  |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed, THEN add all arcs to that variable to the examination queue. ${ }_{66}$

## Constraint Propagation Example AC-3

Graph Coloring
Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $V_{1}-V_{2}$ | none |
| $V_{1}-V_{3}$ | $V_{1}(G)$ |
| $V_{2}-V_{3}$ | $V_{2}(G)$ |
| $V_{2}-V_{1}$ | $V_{1}(R)$ |
| $V_{2}>V_{1}$ |  |
|  |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed,
THEN add all arcs to that variable to the examination queue. ${ }_{67}^{67}$

## Constraint Propagation Example AC-3



| Arc examined | Value deleted |
| :---: | :---: |
| $V_{1}-V_{2}$ | none |
| $V_{1}-V_{3}$ | $V_{1}(G)$ |
| $V_{2}-V_{3}$ | $V_{2}(G)$ |
| $V_{2}-V_{1}$ | $V_{1}(R)$ |
| $V_{2}>V_{1}$ | none |
|  |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed, THEN add all arcs to that variable to the examination queue. 68

## Constraint Propagation Example AC-3

Graph Coloring
Initial Domains


| Arc examined | Value deleted |
| :---: | :---: |
| $V_{1}-V_{2}$ | none |
| $V_{1}-V_{3}$ | $V_{1}(G)$ |
| $V_{2}-V_{3}$ | $V_{2}(G)$ |
| $V_{2}-V_{1}$ | $V_{1}(R)$ |
| $V_{2}>V_{1}$ | none |
| $V_{3}>V_{1}$ |  |



- Delete unmentioned tail values

IF An element of a variable's domain is removed,
THEN add all arcs to that variable to the examination queue.
69

## Constraint Propagation Example AC-3



| Arc examined | Value deleted |
| :---: | :---: |
| $V_{1}-V_{2}$ | none |
| $V_{1}-V_{3}$ | $V_{1}(G)$ |
| $V_{2}-V_{3}$ | $V_{2}(G)$ |
| $V_{2}-V_{1}$ | $V_{1}(R)$ |
| $V_{2}>V_{1}$ | none |
| $V_{3}>V_{1}$ | none |



Arcs to examine
$\square$

IF examination queue is empty
THEN arc (pairwise) consistent.

## Next: To Solve CSPs we combine arc consistency and search

1. Arc consistency (Constraint propagation),

- Eliminates values that are shown locally to not be a part of any solution.

2. Search

- Explores consequences of committing to particular assignments.

Methods Incorporating Search:

- Standard Search
- BackTrack Search (BT)
- BT with Forward Checking (FC)
- Dynamic Variable Ordering (DVO)
- Iterative Repair
- Backjumping (BJ)


## Outline

- Interpreting line diagrams
- Constraint satisfaction problem (CSPS) aka constraint programs (CP).
- Solving CSPs
- Case study: Scheduling (appendix)


## Real World Example: Scheduling as a CSP

Choose time of activities:

- Observations by the Hubble telescope.
- Jobs performed on machine tools.
- Classes taken for degree.

Variables are activities


Domains Are possible start times (or "chunks" of time)
Constraints 1. Activities that use the same resource cannot overlap in time, and
2. Prerequisites are satisfied.

## Case Study: Course Scheduling

## Given:

- 32 required courses ( $8.01,8.02, \ldots$. 16.410), and
- 8 terms (Fall 1, Spring 1, . . . . . Spring 4).

Find: a legal schedule.
Constraints •Pre-requisites satisfied,

- Courses offered only during certain terms,
- A limited number of courses can be taken per term (say 4), and
- Avoid time conflicts between courses.

Note, traditional CSPs are not for expressing (soft) preferences e.g. minimize difficulty, balance subject areas, etc.

But see recent research on valued CSPs!

## Alternative formulations for variables and values <br> DOMAINS

## VARIABLES

A. 1 var per Term
(Fall 1) (Spring 1)
(Fall 2) (Spring 2) . . .

All legal combinations of 4 courses, all offered during that term.
B. 1 var per Term-Slot
subdivide each term into 4 course slots:
(Fall 1,1) (Fall 1, 2)
(Fall1, 3) (Fall 1, 4)
C. 1 var per Course

Terms or term-slots.
Term-slots make it easier to express the constraint limiting the number of courses per term.


MIT OpenCourseWare
http://ocw.mit.edu

### 16.410 / 16.413 Principles of Autonomy and Decision Making

Fall 2010

For information about citing these materials or our Terms of Use, visit: http://ocw.mit.edu/terms.

